

Markov Chains

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September 23, 2020

Definition and examples

Chapman-Kolmogorov equations

Gambler's ruin problem

Queues in communication networks: Transition probabilities

Classes of states

- ▶ Consider discrete-time index $n = 0, 1, 2, \dots$
- ▶ Time-dependent random state X_n takes values on a countable set
 - ▶ In general, states are $i = 0, \pm 1, \pm 2, \dots$, i.e., here the **state space** is \mathbb{Z}
 - ▶ If $X_n = i$ we say “the process is in state i at time n ”
- ▶ Random process is $X_{\mathbb{N}}$, its history up to n is $\mathbf{X}_n = [X_n, X_{n-1}, \dots, X_0]^T$
- ▶ **Def:** process $X_{\mathbb{N}}$ is a **Markov chain (MC)** if for all $n \geq 1, i, j, \mathbf{x} \in \mathbb{Z}^n$
$$P(X_{n+1} = j \mid X_n = i, \mathbf{X}_{n-1} = \mathbf{x}) = P(X_{n+1} = j \mid X_n = i) = P_{ij}$$
- ▶ Future depends only on current state X_n (**memoryless, Markov property**)
 \Rightarrow Future conditionally independent of the past, given the present

- ▶ Given X_n , history \mathbf{X}_{n-1} irrelevant for future evolution of the process
- ▶ From the Markov property, can show that for arbitrary $m > 0$

$$P(X_{n+m} = j \mid X_n = i, \mathbf{X}_{n-1} = \mathbf{x}) = P(X_{n+m} = j \mid X_n = i)$$

- ▶ **Transition probabilities** P_{ij} are constant (MC is time invariant)

$$P(X_{n+1} = j \mid X_n = i) = P(X_1 = j \mid X_0 = i) = P_{ij}$$

- ▶ Since P_{ij} 's are probabilities they are non-negative and sum up to 1

$$P_{ij} \geq 0, \quad \sum_{j=0}^{\infty} P_{ij} = 1$$

⇒ **Conditional probabilities satisfy the axioms**

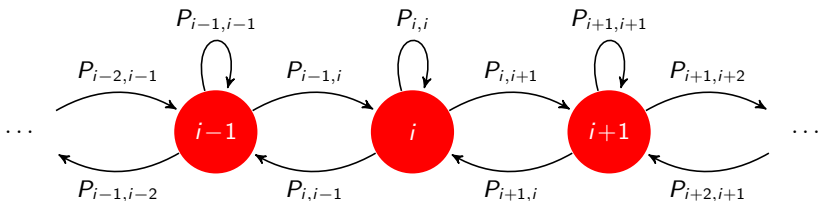
- ▶ Group the P_{ij} in a **transition probability** “matrix” \mathbf{P}

$$\mathbf{P} = \begin{pmatrix} P_{00} & P_{01} & P_{02} & \dots & P_{0j} & \dots \\ P_{10} & P_{11} & P_{12} & \dots & P_{1j} & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ P_{i0} & P_{i1} & P_{i2} & \dots & P_{ij} & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \ddots \end{pmatrix}$$

⇒ Not really a matrix if number of states is infinite

- ▶ **Row-wise** sums should be equal to one, i.e., $\sum_{j=0}^{\infty} P_{ij} = 1$ for all i

- ▶ A graph representation or **state transition diagram** is also used

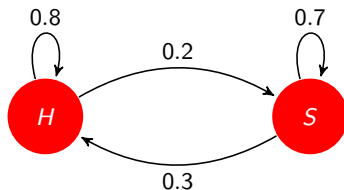


- ▶ Useful when number of states is infinite, skip arrows if $P_{ij} = 0$
- ▶ Again, sum of per-state **outgoing** arrow weights should be one

Example: Happy - Sad

- ▶ I can be happy ($X_n = 0$) or sad ($X_n = 1$)
 \Rightarrow My mood tomorrow is only affected by my mood today
- ▶ Model as Markov chain with transition probabilities

$$\mathbf{P} = \begin{pmatrix} 0.8 & 0.2 \\ 0.3 & 0.7 \end{pmatrix}$$

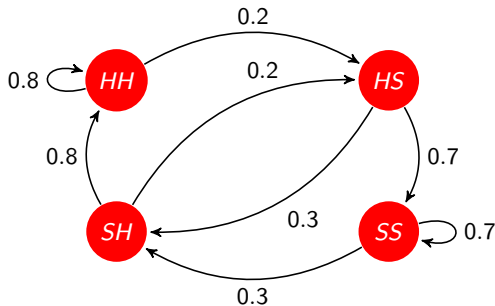


- ▶ Inertia \Rightarrow happy or sad today, likely to stay happy or sad tomorrow
- ▶ But when sad, a little less likely so ($P_{00} > P_{11}$)

Example: Happy - Sad with memory

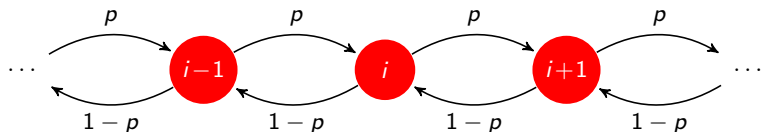
- ▶ Happiness tomorrow affected by today's and yesterday's mood
 - ⇒ Not a Markov chain with the previous state space
- ▶ Define double states HH (Happy-Happy), HS (Happy-Sad), SH, SS
- ▶ Only some transitions are possible
 - ▶ HH and SH can only become HH or HS
 - ▶ HS and SS can only become SH or SS

$$\mathbf{P} = \begin{pmatrix} 0.8 & 0.2 & 0 & 0 \\ 0 & 0 & 0.3 & 0.7 \\ 0.8 & 0.2 & 0 & 0 \\ 0 & 0 & 0.3 & 0.7 \end{pmatrix}$$



- ▶ **Key:** can capture longer time memory via state augmentation

- ▶ Step to the right w.p. p , to the left w.p. $1 - p$
⇒ Not that drunk to stay on the same place

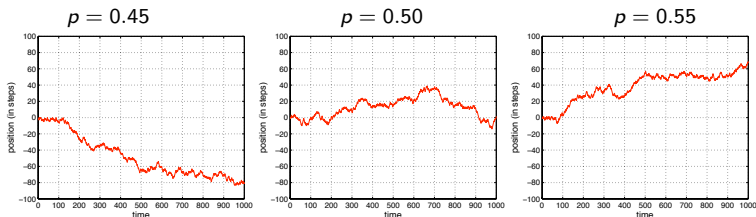


- ▶ States are $0, \pm 1, \pm 2, \dots$ (state space is \mathbb{Z}), **infinite number of states**
- ▶ Transition probabilities are

$$P_{i,i+1} = p, \quad P_{i,i-1} = 1 - p$$

- ▶ $P_{ij} = 0$ for all other transitions

- ▶ Random walks behave differently if $p < 1/2$, $p = 1/2$ or $p > 1/2$



- ⇒ With $p > 1/2$ diverges to the right (\nearrow almost surely)
 - ⇒ With $p < 1/2$ diverges to the left (\searrow almost surely)
 - ⇒ With $p = 1/2$ always come back to visit origin (almost surely)
- ▶ Because number of states is infinite we can have all states transient
 - ▶ **Transient states** not revisited after some time (more later)

Two dimensional random walk

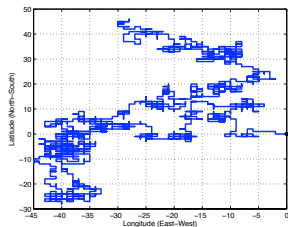
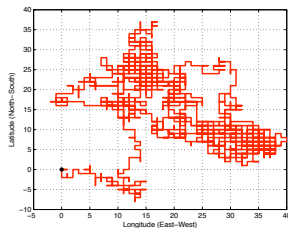
- ▶ Take a step in random direction E, W, S or N
⇒ E, W, S, N chosen with equal probability
- ▶ States are pairs of coordinates (X_n, Y_n)
 - ▶ $X_n = 0, \pm 1, \pm 2, \dots$ and $Y_n = 0, \pm 1, \pm 2, \dots$
- ▶ Transition probs. $\neq 0$ only for adjacent points

$$\text{East: } P(X_{n+1} = i+1, Y_{n+1} = j \mid X_n = i, Y_n = j) = \frac{1}{4}$$

$$\text{West: } P(X_{n+1} = i-1, Y_{n+1} = j \mid X_n = i, Y_n = j) = \frac{1}{4}$$

$$\text{North: } P(X_{n+1} = i, Y_{n+1} = j+1 \mid X_n = i, Y_n = j) = \frac{1}{4}$$

$$\text{South: } P(X_{n+1} = i, Y_{n+1} = j-1 \mid X_n = i, Y_n = j) = \frac{1}{4}$$



- ▶ Some random facts of life for **equiprobable** random walks
- ▶ In one and two dimensions probability of returning to origin is 1
 - ⇒ Will almost surely return home
- ▶ In more than two dimensions, probability of returning to origin is < 1
 - ⇒ In three dimensions probability of returning to origin is 0.34
 - ⇒ Then 0.19, 0.14, 0.10, 0.08, ...

- ▶ Consider an i.i.d. sequence of RVs $Y_{\mathbb{N}} = Y_1, Y_2, \dots, Y_n, \dots$
- ▶ Y_n takes the value ± 1 , $P(Y_n = 1) = p$, $P(Y_n = -1) = 1 - p$
- ▶ Define $X_0 = 0$ and the cumulative sum

$$X_n = \sum_{k=1}^n Y_k$$

\Rightarrow The process $X_{\mathbb{N}}$ is a **random walk** (same we saw earlier)

$\Rightarrow Y_{\mathbb{N}}$ are i.i.d. **steps** (increments) because $X_n = X_{n-1} + Y_n$

- ▶ **Q:** Can we formally establish the random walk is a Markov chain?
- ▶ **A:** Since $X_n = X_{n-1} + Y_n$, $n \geq 1$, and Y_n independent of \mathbf{X}_{n-1}

$$\begin{aligned} P(X_n = j \mid X_{n-1} = i, \mathbf{X}_{n-2} = \mathbf{x}) &= P(X_{n-1} + Y_n = j \mid X_{n-1} = i, \mathbf{X}_{n-2} = \mathbf{x}) \\ &= P(Y_1 = j - i) := P_{ij} \end{aligned}$$

Theorem

Suppose $Y_{\mathbb{N}} = Y_1, Y_2, \dots, Y_n, \dots$ are i.i.d. and independent of X_0 . Consider the random process $X_{\mathbb{N}} = X_1, X_2, \dots, X_n, \dots$ of the form

$$X_n = f(X_{n-1}, Y_n), \quad n \geq 1$$

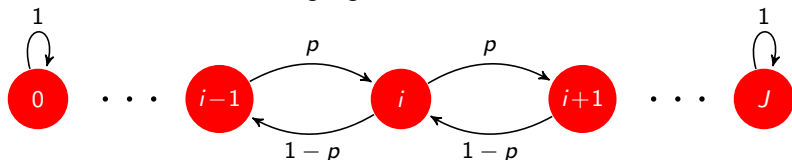
Then $X_{\mathbb{N}}$ is a Markov chain with transition probabilities

$$P_{ij} = P(f(i, Y_1) = j)$$

- ▶ Useful result to identify Markov chains
 - ⇒ Often simpler than checking the Markov property
- ▶ Proof similar to the random walk special case, i.e., $f(x, y) = x + y$

Random walk with boundaries (gambling)

- ▶ As a random walk, but stop moving when $X_n = 0$ or $X_n = J$
 - ▶ Models a gambler that stops playing when ruined, $X_n = 0$
 - ▶ Or when reaches target gains $X_n = J$



- ▶ States are $0, 1, \dots, J$, **finite number of states**
- ▶ Transition probabilities are

$$P_{i,i+1} = p, \quad P_{i,i-1} = 1 - p, \quad P_{00} = 1, \quad P_{JJ} = 1$$

- ▶ $P_{ij} = 0$ for all other transitions
- ▶ States 0 and J are called **absorbing**. Once there stay there forever
⇒ The rest are **transient states**. Visits stop almost surely

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Queues in communication networks: Transition probabilities

Classes of states

- ▶ **Q:** What can be said about multiple transitions?
- ▶ **Ex:** Transition probabilities between two time slots

$$P_{ij}^2 = P(X_{m+2} = j \mid X_m = i)$$

⇒ **Caution:** P_{ij}^2 is just notation, $P_{ij}^2 \neq P_{ij} \times P_{ij}$

- ▶ **Ex:** Probabilities of X_{m+n} given X_m ⇒ **n -step transition probabilities**

$$P_{ij}^n = P(X_{m+n} = j \mid X_m = i)$$

- ▶ Relation between n -, m -, and $(m+n)$ -step transition probabilities
⇒ Write P_{ij}^{m+n} in terms of P_{ij}^m and P_{ij}^n
- ▶ All questions answered by Chapman-Kolmogorov's equations

- ▶ Start considering transition probabilities between two time slots

$$P_{ij}^2 = P(X_{n+2} = j \mid X_n = i)$$

- ▶ Using the law of total probability

$$P_{ij}^2 = \sum_{k=0}^{\infty} P(X_{n+2} = j \mid X_{n+1} = k, X_n = i) P(X_{n+1} = k \mid X_n = i)$$

- ▶ In the first probability, conditioning on $X_n = i$ is unnecessary. Thus

$$P_{ij}^2 = \sum_{k=0}^{\infty} P(X_{n+2} = j \mid X_{n+1} = k) P(X_{n+1} = k \mid X_n = i)$$

- ▶ Which by definition of transition probabilities yields

$$P_{ij}^2 = \sum_{k=0}^{\infty} P_{kj} P_{ik}$$

Relating n -, m -, and $(m + n)$ -step probabilities

- ▶ Same argument works (condition on X_0 w.l.o.g., time invariance)

$$P_{ij}^{m+n} = P(X_{n+m} = j \mid X_0 = i)$$

- ▶ Use law of total probability, drop unnecessary conditioning and use definitions of n -step and m -step transition probabilities

$$P_{ij}^{m+n} = \sum_{k=0}^{\infty} P(X_{m+n} = j \mid X_m = k, X_0 = i) P(X_m = k \mid X_0 = i)$$

$$P_{ij}^{m+n} = \sum_{k=0}^{\infty} P(X_{m+n} = j \mid X_m = k) P(X_m = k \mid X_0 = i)$$

$$P_{ij}^{m+n} = \sum_{k=0}^{\infty} P_{kj}^n P_{ik}^m \quad \text{for all } i, j \text{ and } n, m \geq 0$$

⇒ These are the Chapman-Kolmogorov equations

- ▶ Chapman-Kolmogorov equations are intuitive. Recall

$$P_{ij}^{m+n} = \sum_{k=0}^{\infty} P_{ik}^m P_{kj}^n$$

- ▶ Between times 0 and $m+n$, time m occurred
- ▶ At time m , the Markov chain is in some state $X_m = k$
 - ⇒ P_{ik}^m is the probability of going from $X_0 = i$ to $X_m = k$
 - ⇒ P_{kj}^n is the probability of going from $X_m = k$ to $X_{m+n} = j$
 - ⇒ Product $P_{ik}^m P_{kj}^n$ is then the probability of going from $X_0 = i$ to $X_{m+n} = j$ passing through $X_m = k$ at time m
- ▶ Since any k might have occurred, just sum over all k

- ▶ Define the following three matrices:
 - ⇒ $\mathbf{P}^{(m)}$ with elements P_{ij}^m
 - ⇒ $\mathbf{P}^{(n)}$ with elements P_{ij}^n
 - ⇒ $\mathbf{P}^{(m+n)}$ with elements P_{ij}^{m+n}
- ▶ Matrix product $\mathbf{P}^{(m)}\mathbf{P}^{(n)}$ has (i, j) -th element $\sum_{k=0}^{\infty} P_{ik}^m P_{kj}^n$
- ▶ Chapman Kolmogorov in matrix form

$$\mathbf{P}^{(m+n)} = \mathbf{P}^{(m)}\mathbf{P}^{(n)}$$

- ▶ Matrix of $(m + n)$ -step transitions is product of m -step and n -step

- ▶ For $m = n = 1$ (2-step transition probabilities) matrix form is

$$\mathbf{P}^{(2)} = \mathbf{P}\mathbf{P} = \mathbf{P}^2$$

- ▶ Proceed recursively backwards from n

$$\mathbf{P}^{(n)} = \mathbf{P}^{(n-1)}\mathbf{P} = \mathbf{P}^{(n-2)}\mathbf{P}\mathbf{P} = \dots = \mathbf{P}^n$$

- ▶ Have proved the following

Theorem

The matrix of n -step transition probabilities $\mathbf{P}^{(n)}$ is given by the n -th power of the transition probability matrix \mathbf{P} , i.e.,

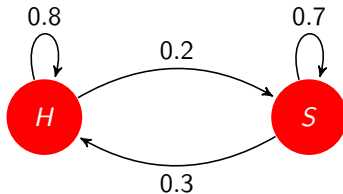
$$\mathbf{P}^{(n)} = \mathbf{P}^n$$

Henceforth we write \mathbf{P}^n

Example: Happy-Sad

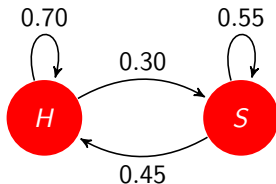
- ▶ Mood transitions in one day

$$\mathbf{P} = \begin{pmatrix} 0.8 & 0.2 \\ 0.3 & 0.7 \end{pmatrix}$$



- ▶ Transition probabilities between today and the day after tomorrow?

$$\mathbf{P}^2 = \begin{pmatrix} 0.70 & 0.30 \\ 0.45 & 0.55 \end{pmatrix}$$



- ▶ ... After a week and after a month

$$\mathbf{P}^7 = \begin{pmatrix} 0.6031 & 0.3969 \\ 0.5953 & 0.4047 \end{pmatrix} \quad \mathbf{P}^{30} = \begin{pmatrix} 0.6000 & 0.4000 \\ 0.6000 & 0.4000 \end{pmatrix}$$

- ▶ Matrices \mathbf{P}^7 and \mathbf{P}^{30} almost identical $\Rightarrow \lim_{n \rightarrow \infty} \mathbf{P}^n$ exists
 \Rightarrow Note that this is a regular limit
- ▶ After a month transition from H to H and from S to H w.p. 0.6
 \Rightarrow State becomes independent of initial condition (H w.p. 0.6)
- ▶ **Rationale:** 1-step memory \Rightarrow Initial condition eventually forgotten
 - ▶ More about this soon

- ▶ All probabilities so far are conditional, i.e., $P_{ij}^n = P(X_n = j \mid X_0 = i)$
⇒ May want **unconditional probabilities** $p_j(n) = P(X_n = j)$
- ▶ Requires specification of **initial conditions** $p_i(0) = P(X_0 = i)$
- ▶ Using law of total probability and definitions of P_{ij}^n and $p_j(n)$

$$\begin{aligned} p_j(n) = P(X_n = j) &= \sum_{i=0}^{\infty} P(X_n = j \mid X_0 = i) P(X_0 = i) \\ &= \sum_{i=0}^{\infty} P_{ij}^n p_i(0) \end{aligned}$$

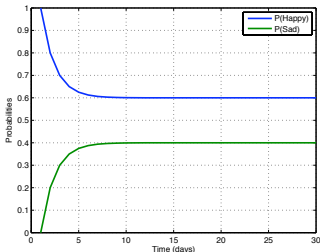
- ▶ In matrix form (define vector $\mathbf{p}(n) = [p_1(n), p_2(n), \dots]^T$)

$$\mathbf{p}(n) = (\mathbf{P}^n)^T \mathbf{p}(0)$$

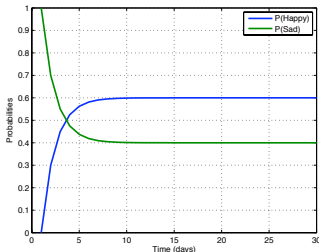
Example: Happy-Sad

► Transition probability matrix $\Rightarrow \mathbf{P} = \begin{pmatrix} 0.8 & 0.2 \\ 0.3 & 0.7 \end{pmatrix}$

$$\mathbf{p}(0) = [1, 0]^T$$



$$\mathbf{p}(0) = [0, 1]^T$$



► For large n probabilities $\mathbf{p}(n)$ are independent of initial state $\mathbf{p}(0)$

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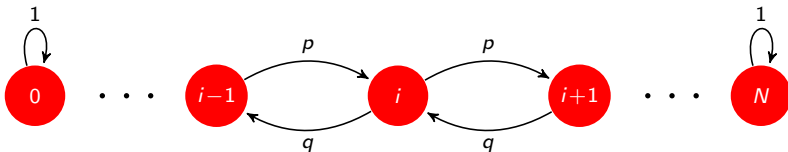
Queues in communication networks: Transition probabilities

Classes of states

- ▶ You place \$1 bets
 - (i) With probability p you gain \$1, and
 - (ii) With probability $q = 1 - p$ you lose your \$1 bet
- ▶ Start with an initial wealth of $\$i$
- ▶ Define bias factor $\alpha := q/p$
 - ▶ If $\alpha > 1$ more likely to lose than win (biased against gambler)
 - ▶ $\alpha < 1$ favors gambler (more likely to win than lose)
 - ▶ $\alpha = 1$ game is fair
- ▶ You keep playing until
 - (a) You go broke (lose all your money)
 - (b) You reach a wealth of $\$N$ (same as first lecture, HW1 for $N \rightarrow \infty$)
- ▶ Prob. S_i of reaching $\$N$ before going broke for initial wealth $\$i$?
 - ▶ S stands for success, or successful betting run (SBR)

- ▶ Model wealth as **Markov chain** X_N . Transition probabilities

$$P_{i,i+1} = p, \quad P_{i,i-1} = q, \quad P_{00} = P_{NN} = 1$$



- ▶ **Realizations** x_N . Initial state = Initial wealth = i
 - ⇒ States 0 and N are **absorbing**. Eventually end up in one of them
 - ⇒ Remaining states are **transient** (visits eventually stop)
- ▶ Being absorbing states says something about the **limit wealth**

$$\lim_{n \rightarrow \infty} x_n = 0, \text{ or } \lim_{n \rightarrow \infty} x_n = N \Rightarrow S_i := P \left(\lim_{n \rightarrow \infty} X_n = N \mid X_0 = i \right)$$

- ▶ Total probability to relate S_i with S_{i+1}, S_{i-1} from adjacent states
⇒ Condition on first bet X_1 , Markov chain homogeneous

$$S_i = S_{i+1}P_{i,i+1} + S_{i-1}P_{i,i-1} = S_{i+1}p + S_{i-1}q$$

- ▶ Recall $p + q = 1$ and reorder terms

$$p(S_{i+1} - S_i) = q(S_i - S_{i-1})$$

- ▶ Recall definition of bias $\alpha = q/p$

$$S_{i+1} - S_i = \alpha(S_i - S_{i-1})$$

- ▶ If current state is 0 then $S_i = S_0 = 0$. Can write

$$S_2 - S_1 = \alpha(S_1 - S_0) = \alpha S_1$$

- ▶ Substitute this in the expression for $S_3 - S_2$

$$S_3 - S_2 = \alpha(S_2 - S_1) = \alpha^2 S_1$$

- ▶ Apply recursively backwards from $S_i - S_{i-1}$

$$S_i - S_{i-1} = \alpha(S_{i-1} - S_{i-2}) = \dots = \alpha^{i-1} S_1$$

- ▶ Sum up all of the former to obtain

$$S_i - S_1 = S_1(\alpha + \alpha^2 + \dots + \alpha^{i-1})$$

- ▶ The latter can be written as a geometric series

$$S_i = S_1(1 + \alpha + \alpha^2 + \dots + \alpha^{i-1})$$

- ▶ Geometric series can be summed in closed form, assuming $\alpha \neq 1$

$$S_i = \left(\sum_{k=0}^{i-1} \alpha^k \right) S_1 = \frac{1 - \alpha^i}{1 - \alpha} S_1$$

- ▶ When in state N , $S_N = 1$ and so

$$1 = S_N = \frac{1 - \alpha^N}{1 - \alpha} S_1 \Rightarrow S_1 = \frac{1 - \alpha}{1 - \alpha^N}$$

- ▶ Substitute S_1 above into expression for probability of SBR S_i

$$S_i = \frac{1 - \alpha^i}{1 - \alpha^N}, \quad \alpha \neq 1$$

- ▶ For $\alpha = 1 \Rightarrow S_i = iS_1$, $1 = S_N = NS_1$, $\Rightarrow S_i = \frac{i}{N}$

- ▶ Recall

$$S_i = \begin{cases} (1 - \alpha^i)/(1 - \alpha^N), & \alpha \neq 1, \\ i/N, & \alpha = 1 \end{cases}$$

- ▶ Consider exit bound N arbitrarily large

(i) For $\alpha > 1$, $S_i \approx (\alpha^i - 1)/\alpha^N \rightarrow 0$

(ii) Likewise for $\alpha = 1$, $S_i = i/N \rightarrow 0$

- ▶ If win probability p does not exceed loose probability q
 \Rightarrow Will almost surely lose all money

(iii) For $\alpha < 1$, $S_i \rightarrow 1 - \alpha^i$

- ▶ If win probability p exceeds loose probability q
 \Rightarrow For sufficiently high initial wealth i , will most likely win
- ▶ This explains what we saw on first lecture and HW1

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Queues in communication networks: Transition probabilities

Classes of states

- ▶ General **communication systems** goal
 - ⇒ Move packets from generating sources to intended destinations
- ▶ Between arrival and departure we hold packets in a memory buffer
 - ⇒ **Want to design buffers appropriately**

- ▶ Time slotted in intervals of duration Δt
 - ⇒ n -th slot between times $n\Delta t$ and $(n + 1)\Delta t$
- ▶ Average arrival rate is $\bar{\lambda}$ packets per unit time
 - ⇒ Probability of packet arrival in Δt is $\lambda = \bar{\lambda}\Delta t$
- ▶ Packets are transmitted (depart) at a rate of $\bar{\mu}$ packets per unit time
 - ⇒ Probability of packet departure in Δt is $\mu = \bar{\mu}\Delta t$
- ▶ Assume no simultaneous arrival and departure (no concurrence)
 - ⇒ Reasonable for small Δt (μ and λ likely to be small)

- ▶ Q_n denotes number of packets in queue (backlog) in n -th time slot
- ▶ $\mathbb{A}_n =$ nr. of packet arrivals, $\mathbb{D}_n =$ nr. of departures (during n -th slot)
- ▶ If the queue is empty $Q_n = 0$ then there are no departures
⇒ Queue length at time $n + 1$ can be written as

$$Q_{n+1} = Q_n + \mathbb{A}_n, \quad \text{if } Q_n = 0$$

- ▶ If $Q_n > 0$, departures and arrivals may happen

$$Q_{n+1} = Q_n + \mathbb{A}_n - \mathbb{D}_n, \quad \text{if } Q_n > 0$$

- ▶ $\mathbb{A}_n \in \{0, 1\}$, $\mathbb{D}_n \in \{0, 1\}$ and either $\mathbb{A}_n = 1$ or $\mathbb{D}_n = 1$ but not both
⇒ Arrival and departure probabilities are

$$P(\mathbb{A}_n = 1) = \lambda, \quad P(\mathbb{D}_n = 1) = \mu$$

- ▶ Future queue lengths depend on current length only
- ▶ Probability of queue length increasing

$$P(Q_{n+1} = i + 1 \mid Q_n = i) = P(\mathbb{A}_n = 1) = \lambda, \quad \text{for all } i$$

- ▶ Queue length might decrease only if $Q_n > 0$. Probability is

$$P(Q_{n+1} = i - 1 \mid Q_n = i) = P(\mathbb{D}_n = 1) = \mu, \quad \text{for all } i > 0$$

- ▶ Queue length stays the same if it neither increases nor decreases

$$P(Q_{n+1} = i \mid Q_n = i) = 1 - \lambda - \mu, \quad \text{for all } i > 0$$

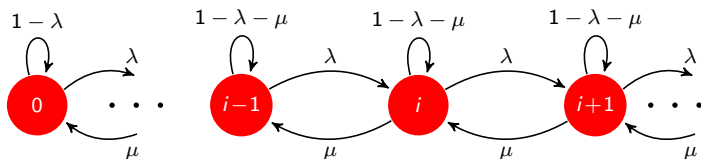
$$P(Q_{n+1} = 0 \mid Q_n = 0) = 1 - \lambda$$

⇒ No departures when $Q_n = 0$ explain second equation

- ▶ MC with states $0, 1, 2, \dots$. Identify states with queue lengths
- ▶ Transition probabilities for $i \neq 0$ are

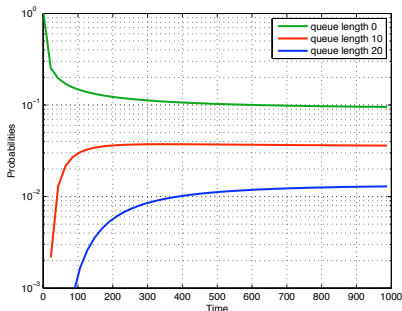
$$P_{i,i-1} = \mu, \quad P_{i,i} = 1 - \lambda - \mu, \quad P_{i,i+1} = \lambda$$

- ▶ For $i = 0$: $P_{00} = 1 - \lambda$ and $P_{01} = \lambda$



Numerical example: Probability propagation

- ▶ Build matrix \mathbf{P} truncating at maximum queue length $L = 100$
 - ⇒ Arrival rate $\lambda = 0.3$. Departure rate $\mu = 0.33$
 - ⇒ Initial distribution $\mathbf{p}(0) = [1, 0, 0, \dots]^T$ (queue empty)



- ▶ Propagate probabilities $(\mathbf{P}^n)^T \mathbf{p}(0)$
- ▶ Probabilities obtained are

$$P(Q_n = i \mid Q_0 = 0) = p_i(n)$$

- ▶ A few i 's (0, 10, 20) shown
- ▶ Probability of empty queue ≈ 0.1
- ▶ Occupancy decreases with i

Definition and examples

Chapman-Kolmogorov equations

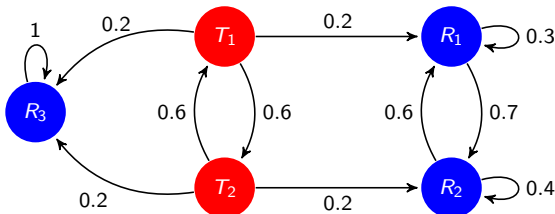
Gambler's ruin problem

Queues in communication networks: Transition probabilities

Classes of states

Transient and recurrent states

- ▶ States of a MC can be **recurrent** or **transient**
- ▶ **Transient states** might be visited early on but visits eventually stop
- ▶ Almost surely, $X_n \neq i$ for n sufficiently large (qualifications needed)
- ▶ Visits to **recurrent states** keep happening forever. Fix arbitrary m
- ▶ Almost surely, $X_n = i$ for some $n \geq m$ (qualifications needed)



- ▶ Let f_i be the probability that starting at i , MC ever reenters state i

$$f_i := P\left(\bigcup_{n=1}^{\infty} X_n = i \mid X_0 = i\right) = P\left(\bigcup_{n=m+1}^{\infty} X_n = i \mid X_m = i\right)$$

- ▶ State i is **recurrent** if $f_i = 1$
 - ⇒ Process reenters i again and again (a.s.). **Infinitely often**
- ▶ State i is **transient** if $f_i < 1$
 - ⇒ Positive probability $1 - f_i > 0$ of never coming back to i

- ▶ State R_3 is **recurrent** because it is absorbing $P(X_1 = R_3 | X_0 = R_3) = 1$

- ▶ State R_1 is **recurrent** because

$$P(X_1 = R_1 | X_0 = R_1) = 0.3$$

$$P(X_2 = R_1, X_1 \neq R_1 | X_0 = R_1) = (0.7)(0.6)$$

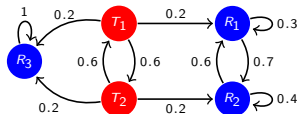
$$P(X_3 = R_1, X_2 \neq R_1, X_1 \neq R_1 | X_0 = R_1) = (0.7)(0.4)(0.6)$$

⋮

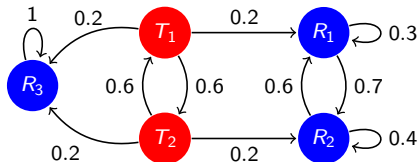
$$P(X_n = R_1, X_{n-1} \neq R_1, \dots, X_1 \neq R_1 | X_0 = R_1) = (0.7)(0.4)^{n-2}(0.6)$$

- ▶ Sum up: $f_i = \sum_{n=1}^{\infty} P(X_n = R_1, X_{n-1} \neq R_1, \dots, X_1 \neq R_1 | X_0 = R_1)$

$$= 0.3 + 0.7 \left(\sum_{n=2}^{\infty} 0.4^{n-2} \right) 0.6 = 0.3 + 0.7 \left(\frac{1}{1 - 0.4} \right) 0.6 = 1$$



- ▶ States T_1 and T_2 are **transient**
- ▶ Probability of returning to T_1 is $f_{T_1} = (0.6)^2 = 0.36$
 - ⇒ Might come back to T_1 only if it goes to T_2 (w.p. 0.6)
 - ⇒ Will come back only if it moves back from T_2 to T_1 (w.p. 0.6)



- ▶ Likewise, $f_{T_2} = (0.6)^2 = 0.36$

- ▶ Define N_i as the number of visits to state i given that $X_0 = i$

$$N_i := \sum_{n=1}^{\infty} \mathbb{I} \{X_n = i \mid X_0 = i\}$$

- ▶ If $X_n = i$, this is the last visit to i w.p. $1 - f_i$
- ▶ Prob. revisiting state i exactly n times is (n visits \times no more visits)

$$P(N_i = n) = f_i^n(1 - f_i)$$

\Rightarrow Number of visits $N_i + 1$ is geometric with parameter $1 - f_i$

- ▶ Expected number of visits is

$$\mathbb{E}[N_i] + 1 = \frac{1}{1 - f_i} \Rightarrow \mathbb{E}[N_i] = \frac{f_i}{1 - f_i}$$

\Rightarrow For recurrent states $N_i = \infty$ a.s. and $\mathbb{E}[N_i] = \infty$ ($f_i = 1$)

- ▶ Another way of writing $\mathbb{E}[N_i]$

$$\mathbb{E}[N_i] = \sum_{n=1}^{\infty} \mathbb{E}[\mathbb{I}\{X_n = i \mid X_0 = i\}] = \sum_{n=1}^{\infty} P_{ii}^n$$

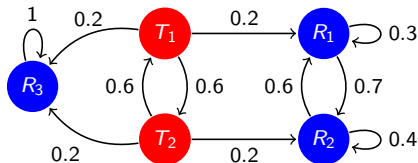
- ▶ Recall that: for **transient** states $\mathbb{E}[N_i] = f_i/(1 - f_i) < \infty$
for **recurrent** states $\mathbb{E}[N_i] = \infty$

Theorem

- ▶ State i is **transient** if and only if $\sum_{n=1}^{\infty} P_{ii}^n < \infty$
- ▶ State i is **recurrent** if and only if $\sum_{n=1}^{\infty} P_{ii}^n = \infty$

- ▶ Number of future visits to **transient** states is **finite**
⇒ If number of states is **finite** some states have to be **recurrent**

- ▶ **Def:** State j is **accessible** from state i if $P_{ij}^n > 0$ for some $n \geq 0$
⇒ It is possible to enter j if MC initialized at $X_0 = i$
- ▶ Since $P_{ii}^0 = P(X_0 = i | X_0 = i) = 1$, **state i is accessible from itself**



- ▶ All states accessible from T_1 and T_2
- ▶ Only R_1 and R_2 accessible from R_1 or R_2
- ▶ None other than R_3 accessible from itself

- ▶ **Def:** States i and j are said to **communicate** ($i \leftrightarrow j$) if
 - ⇒ j is accessible from i , i.e., $P_{ij}^n > 0$ for some n ; and
 - ⇒ i is accessible from j , i.e., $P_{ji}^m > 0$ for some m
- ▶ **Communication is an equivalence relation**
- ▶ **Reflexivity:** $i \leftrightarrow i$
 - ▶ Holds because $P_{ii}^0 = 1$
- ▶ **Symmetry:** If $i \leftrightarrow j$ then $j \leftrightarrow i$
 - ▶ If $i \leftrightarrow j$ then $P_{ij}^n > 0$ and $P_{ji}^m > 0$ from where $j \leftrightarrow i$
- ▶ **Transitivity:** If $i \leftrightarrow j$ and $j \leftrightarrow k$, then $i \leftrightarrow k$
 - ▶ Just notice that $P_{ik}^{n+m} \geq P_{ij}^n P_{jk}^m > 0$
- ▶ **Partitions set of states into disjoint classes** (as all equivalences do)
 - ⇒ What are these classes?

Theorem

If state i is *recurrent* and $i \leftrightarrow j$, then j is *recurrent*

Proof.

- ▶ If $i \leftrightarrow j$ then there are l, m such that $P_{ji}^l > 0$ and $P_{ij}^m > 0$
- ▶ Then, for any n we have

$$P_{jj}^{l+n+m} \geq P_{ji}^l P_{ii}^n P_{ij}^m$$

- ▶ Sum for all n . Note that since i is recurrent $\sum_{n=1}^{\infty} P_{ii}^n = \infty$

$$\sum_{n=1}^{\infty} P_{jj}^{l+n+m} \geq \sum_{n=1}^{\infty} P_{ji}^l P_{ii}^n P_{ij}^m = P_{ji}^l \left(\sum_{n=1}^{\infty} P_{ii}^n \right) P_{ij}^m = \infty$$

\Rightarrow Which implies j is recurrent



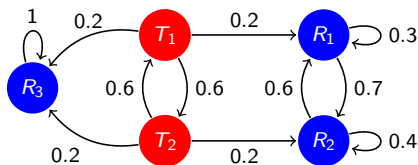
Corollary

If state i is *transient* and $i \leftrightarrow j$, then j is *transient*

Proof.

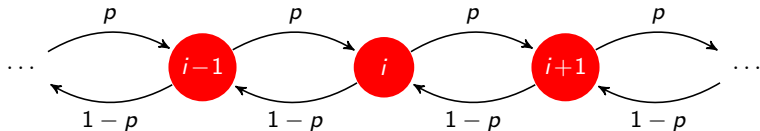
- ▶ If j were recurrent, then i would be recurrent from previous theorem □
- ▶ Recurrence is shared by elements of a communication class
⇒ We say that recurrence is a class property
- ▶ Likewise, transience is also a class property
- ▶ MC states are separated in classes of transient and recurrent states

- ▶ A MC is called **irreducible** if it has only one class
 - ▶ All states communicate with each other
 - ▶ If MC also has finite number of states the single class is recurrent
 - ▶ If MC infinite, class might be transient
- ▶ When it has multiple classes (not irreducible)
 - ▶ Classes of transient states $\mathcal{T}_1, \mathcal{T}_2, \dots$
 - ▶ Classes of recurrent states $\mathcal{R}_1, \mathcal{R}_2, \dots$
- ▶ If MC initialized in a recurrent class \mathcal{R}_k , stays within the class
- ▶ If MC starts in transient class \mathcal{T}_k , then it might
 - Stay on \mathcal{T}_k (only if $|\mathcal{T}_k| = \infty$)
 - End up in another transient class \mathcal{T}_r (only if $|\mathcal{T}_r| = \infty$)
 - End up in a recurrent class \mathcal{R}_l
- ▶ For large time index n , MC restricted to one class
 - ⇒ Can be separated into irreducible components



- ▶ Three classes
 - ⇒ $\mathcal{T} := \{T_1, T_2\}$, class with **transient** states
 - ⇒ $\mathcal{R}_1 := \{R_1, R_2\}$, class with **recurrent** states
 - ⇒ $\mathcal{R}_2 := \{R_3\}$, class with **recurrent** state
- ▶ For large n suffices to study the irreducible components \mathcal{R}_1 and \mathcal{R}_2

- ▶ Step right with probability p , left with probability $q = 1 - p$



- ▶ All states communicate \Rightarrow States either all transient or all recurrent
- ▶ To see which, consider initially $X_0 = 0$ and note for any $n \geq 1$

$$P_{00}^{2n} = \binom{2n}{n} p^n q^n = \frac{(2n)!}{n!n!} p^n q^n$$

\Rightarrow Back to 0 in $2n$ steps $\Leftrightarrow n$ steps right and n steps left

Example: Random walk (continued)

- ▶ **Stirling's formula** $n! \approx n^n \sqrt{n} e^{-n} \sqrt{2\pi}$
 ⇒ Approximate probability P_{00}^{2n} of returning home as

$$P_{00}^{2n} = \frac{(2n)!}{n!n!} p^n q^n \approx \frac{(4pq)^n}{\sqrt{n\pi}}$$

- ▶ Symmetric random walk ($p = q = 1/2$)

$$\sum_{n=1}^{\infty} P_{00}^{2n} = \sum_{n=1}^{\infty} \frac{1}{\sqrt{n\pi}} = \infty$$

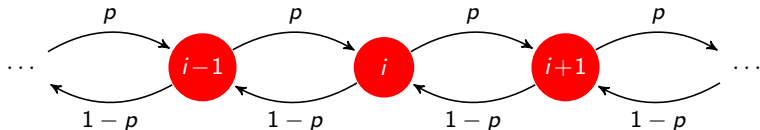
⇒ State 0 (hence all states) are **recurrent**

- ▶ Biased random walk ($p > 1/2$ or $p < 1/2$), then $pq < 1/4$ and

$$\sum_{n=1}^{\infty} P_{00}^{2n} = \sum_{n=1}^{\infty} \frac{(4pq)^n}{\sqrt{n\pi}} < \infty$$

⇒ State 0 (hence all states) are **transient**

- ▶ Alternative proof of **transience** of **right-biased random walk** ($p > 1/2$)



- ▶ Write current position of random walker as $X_n = \sum_{k=1}^n Y_k$
 $\Rightarrow Y_k$ are the i.i.d. steps: $\mathbb{E}[Y_k] = 2p - 1$, $\text{var}[Y_k] = 4p(1-p)$
- ▶ From Central Limit Theorem ($\Phi(x)$ is cdf of standard Normal)

$$P\left(\frac{\sum_{k=1}^n Y_k - n(2p - 1)}{\sqrt{n4p(1-p)}} \leq a\right) \rightarrow \Phi(a)$$

- ▶ Choose $a = \frac{\sqrt{n}(1-2p)}{\sqrt{4p(1-p)}} < 0$, use **Chernoff bound** $\Phi(a) \leq \exp(-a^2/2)$

$$P(X_n \leq 0) = P\left(\sum_{k=1}^n Y_k \leq 0\right) \rightarrow \Phi\left(\frac{\sqrt{n}(1-2p)}{\sqrt{4p(1-p)}}\right) < e^{-\frac{n(1-2p)^2}{8p(1-p)}} \rightarrow 0$$

- ▶ Since $P_{00}^n \leq P(X_n \leq 0)$, sum over n

$$\sum_{n=1}^{\infty} P_{00}^n \leq \sum_{n=1}^{\infty} P(X_n \leq 0) < \sum_{n=1}^{\infty} e^{-\frac{n(1-2p)^2}{8p(1-p)}} < \infty$$

- ▶ This establishes state 0 is **transient**
 \Rightarrow Since all states communicate, **all states are transient**

- ▶ States of a MC can be **transient** or **recurrent**
- ▶ A MC can be partitioned into classes of communicating states
 - ⇒ Class members are either all transient or all recurrent
 - ⇒ **Recurrence and transience are class properties**
 - ⇒ A finite MC has at least one **recurrent** class
- ▶ A MC with only one class is **irreducible**
 - ⇒ If reducible it can be separated into irreducible components

- ▶ Markov chain
- ▶ State space
- ▶ Markov property
- ▶ Transition probability matrix
- ▶ State transition diagram
- ▶ State augmentation
- ▶ Random walk
- ▶ n -step transition probabilities
- ▶ Chapman-Kolmogorov eqs.
- ▶ Initial distribution
- ▶ Gambler's ruin problem
- ▶ Communication system
- ▶ Non-concurrent queue
- ▶ Queue evolution model
- ▶ Recurrent and transient states
- ▶ Accessibility
- ▶ Communication
- ▶ Equivalence relation
- ▶ Communication classes
- ▶ Class property
- ▶ Irreducible Markov chain
- ▶ Irreducible components