

## Characterizing Network Cohesion

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#### Local density, clustering coefficient and group centrality

Network connectivity

Assortativity mixing

Case study: Analysis of an epileptic seizure



Many network analytic questions pertain to network cohesion

#### Example

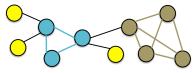
- Q1: Do common friends of an actor end up being friends?
- ▶ Q2: What collections of proteins in a cell work closely together?
- Q3: Does Web page structure separate relative to content?
- Q4: What portion of the Internet topology constitutes a 'backbone'?
- Definitions of network cohesion depend on the context
  - $\Rightarrow$  Scale from local (e.g., triads) to global (e.g., giant components)
  - $\Rightarrow$  Specified explicitly (e.g., cliques) or implicitly (e.g., clusters)



- Cohesive subgroups defined by social network analysts as: 'Actors connected via dense, directed, reciprocated relations'
- Allow sharing information, creating solidarity, collective actions
   Ex: religious cults, terrorist cells, sport clubs, military platoons, ...
- Desirable properties of a cohesive subgroup
  - $\Rightarrow$  Familiarity (degree);
  - $\Rightarrow$  Reachability (distance);
  - $\Rightarrow$  Robustness (connectivity); and
  - $\Rightarrow$  Density (edge density)
- ▶ Natural to think of cliques, i.e., complete subgraphs of G



► Large cliques are rare; single missing edge destroys property



Sufficient condition for the existence of a size-*n* clique

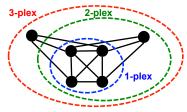
$$N_e > rac{N_v^2}{2} rac{(n-2)}{(n-1)}, ext{ while sparse graphs have } N_e = O(N_v)$$

Complexity of clique-related algorithms varies widely

- ▶ Is  $U \subseteq V$  a clique? Is it maximal?  $O(N_v + N_e)$  complexity
- Identifying all triangles in G?  $O(N_v^3)$  ( $O(N_v^{\sqrt{2}})$  for sparse graphs)
- Does G have a maximal clique of size  $\geq n$ ? NP-complete



- Cliques tend to be an overly restrictive notion of cohesiveness. Relax!
- ▶ Def: An induced subgraph G'(V', E') is a k-plex if d<sub>v</sub>(G') ≥ |V'| k for all v ∈ V', and G' is maximal

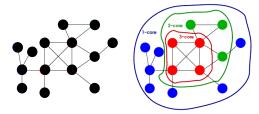


 $\Rightarrow$  Degrees are in the induced subgraph G', not in G

- No vertex is missing more than k − 1 of its possible |V'| − 1 edges ⇒ A clique is a 1-plex
- ► Complex: problems involving *k*-plexes scale like clique counterparts



Recall the k-core decomposition. A dual notion of cohesiveness

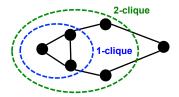


▶ Def: An induced subgraph G'(V', E') is a k-core if d<sub>v</sub>(G') ≥ k for all v ∈ V', and G' is maximal

- $\blacktriangleright$  Hierarchy: larger "coreness"  $\Rightarrow$  larger degrees and centrality
- ► Algorithm: recursively prune all vertices of degree less than k⇒ Complexity  $O(N_v + N_e)$ , very efficient for sparse graphs



- Idea: specify that any two actors are no more than k hops away
- ▶ Def: An induced subgraph G'(V', E') is a k-clique if d(u, v) ≤ k for all u, v ∈ V'



 $\Rightarrow$  Useful if important social processes occur via intermediaries  $\Rightarrow$  diam(G') may exceed k, if distances used are in G

► Likewise, a k-club is a subgraph G' with diam(G') ≤ k ⇒ k-clubs are k-cliques but the converse is not true, in general



• A natural measure of density of a subgraph G'(V', E') is

$$\mathsf{den}(G') = \frac{|E'|}{|V'|(|V'|-1)/2} \in [0,1]$$

 $\Rightarrow$  Quantifies how close is G' to being a clique

• den(G') is just a rescaling of the average degree  $\overline{d}(G')$ 

$$ar{d}(G') = rac{1}{|V'|} \sum_{v \in V'} d_v = rac{2|E'|}{|V'|} \ \Rightarrow \ \mathsf{den}(G') = rac{ar{d}(G')}{|V'|-1}$$

► Flexibility in choosing G' to measure local density via den(G')
 ⇒ Use v's egonet G'<sub>v</sub>, subgraph induced by v and its neighbors
 ⇒ Density of the overall graph G is den(G) = <sup>2Ne</sup>/<sub>N<sub>v</sub>(N<sub>v</sub>-1)</sub>

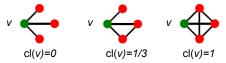
## Clustering coefficient



- ▶ Q: What fraction of *v*'s neighbors are themselves connected?
- **Def:** The clustering coefficient cl(v) of  $v \in V$  is

$$\mathsf{cl}(v) = \frac{2|E_v|}{d_v(d_v-1)} \in [0,1]$$

 $\Rightarrow |E_v|$  is the number of edges among v's neighbors

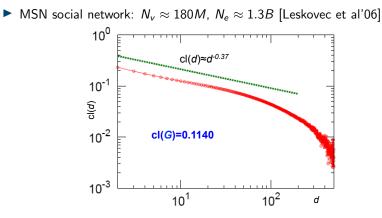


- An indication of the extent to which edges 'cluster'
- The global (average) clustering coefficient is

$$\mathsf{cl}(G) = \frac{1}{N_v} \sum_{v \in V} \mathsf{cl}(v)$$

#### Example: MSN social network





- Average clustering coefficient cl(G) = 0.1140 is large
- Compare with the Erdös-Renyi random graph model

$$\mathsf{cl}(G_{n,p}) = \mathsf{Pr}\left[\mathsf{Edge\ closes\ triangle}
ight] = p = rac{ar{d}}{n-1} o 0$$



- Capture the importance of node subgroups [Everett et al'99]
- ▶ Q1: Are engineers more popular than accountants in an organization?
- Q2: How do we select board members with most business influence?
- Group centrality measures to generalize vertex centrality
- Ex: Consider subgraph G'(V', E') induced by node subset V'
  - Let  $U_{V'} \subset V \setminus V'$  with edges to members of V'
- Group degree centrality of node subset V'

$$d_{V'} = |U_{V'}|$$

 $\Rightarrow$  Number of non-group nodes connected to G'



**Def:** Distance from  $v \in V$  to a group of nodes  $V' \subset V$  is

$$d_*(v,V') = \min_{u \in V'} d(u,v)$$

► Group closeness centrality of node subset V'

$$c_{Cl}(V') = \frac{1}{\sum_{u \in V \setminus V'} d_*(u, V')}$$

► Group betweenness centrality of node subset V'

$$c_{Be}(V') = \sum_{s \neq t \in V \setminus V'} \frac{\sigma(s, t | V')}{\sigma(s, t)}$$

•  $\sigma(s,t)$  is the total number of s-t shortest paths  $(s,t\in V\setminus V')$ 

•  $\sigma(s,t|V')$  is the number of s-t shortest paths through  $v \in V'$ 



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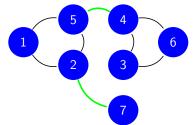
Case study: Analysis of an epileptic seizure



Connectivity relevant when taking a larger, global perspective

- Q: Does a given graph G separate into different subgraphs?
- If it does not, a 'less robust' network is closer to splitting

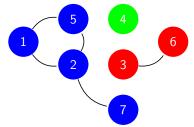
▶ **Def:** Graph is connected if ∃ walks joining each vertex pair



 $\Rightarrow$  If bridge edges are removed, the graph becomes disconnected



A component is a maximally-connected subgraph

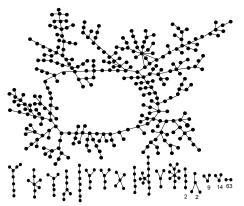


- ▶ In figure  $\Rightarrow$  Components are  $\{1, 2, 5, 7\}$ ,  $\{3, 6\}$  and  $\{4\}$  $\Rightarrow$  Subgraph  $\{3, 4, 6\}$  not connected,  $\{1, 2, 5\}$  not maximal
- Disconnected graphs have 2 or more components
  - $\Rightarrow$  Number of components = Multiplicity of eigenvalue 0 for L
  - $\Rightarrow$  Largest component often called giant component
- Check for connectivity, identify components with DFS, BFS:  $O(N_v)$

### Giant connected components



- Large real-world networks typically exhibit one giant component
- Ex: romantic relationships in a US high school [Bearman et al'04]



Q: Why do we expect to find a single giant component?

A: Well, it only takes one edge to merge two giant components

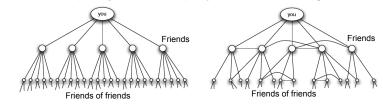
### Average path length and small world



- Giant components tend to exhibit the small world property
- Small refers to the average path length

$$\bar{\ell} = \binom{N_{\nu}}{2}^{-1} \sum_{u \neq v \in V} d(u, v) = O(\log N_{\nu})$$

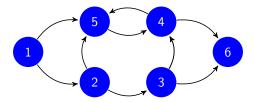
Ex: facilitates spread of gossip, diseases, search for WWW contentNot too surprising that the property holds. Informal argument:



▶ If  $d_v = d$ , after  $h_*$  hops have  $d^{h_*} \approx N_v \implies \bar{\ell} \approx h_* = O(\log N_v)$ 



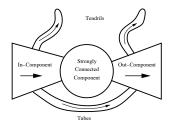
- Connectivity is more subtle with directed graphs. Two notions
- ▶ Def: Digraph is strongly connected if for every pair u, v ∈ V, u is reachable from v (via a directed walk) and vice versa
- Def: Digraph is weakly connected if connected after disregarding arc directions, i.e., the underlying undirected graph is connected



Above graph is weakly connected but not strongly connected
 Strong connectivity obviously implies weak connectivity



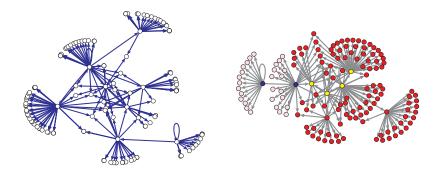
First described for the Web graph in [Broder et al'00]



- Core element is the strongly-connected component (SCC)
  - In-component (IC): vertices reaching SCC, but not vice-versa
  - Out-component (OC): vertices reached by SCC, but not vice-versa
  - Tubes: vertices in between the IC and OC, not in SCC
  - Tendrils: vertices that cannot be reached by, or reach the SCC
- ▶ In general, the digraph may be disconnected with a giant SCC

## Example: AIDS blog network





Network of citations among 146 blogs related to AIDS
 Small SCC with 4 vertices and IC with 2 vertices
 OC dominates with 112 vertices, and few tendrils (28 vertices)

▶ For the WWW, Broder et al. found  $|SCC| \approx |IC| \approx |OC| \approx 56M$ 



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## Assortative mixing



People have a stronger tendency to associate with equals  $\Rightarrow$  Tendency is called homophily or assortative mixing Black White Other

Ex: high-school students by race, bloggers by political party, ...
 ⇒ Can have disassortative mixing e.g., romantic relationships



- Suppose that vertex characteristics are categorical, e.g., male/female
- ► Let  $f_{ij}$  be the fraction of edges joining vertices of categories  $C_i$ ,  $C_j$ ⇒  $f_{i+} = \sum_i f_{ij} (f_{+i})$  is the *i*-th marginal row (column) sum
- Define the assortativity coefficient [Newman'03]

$$r_{a} = \frac{\sum_{i} f_{ii} - \sum_{i} f_{i+} f_{+i}}{1 - \sum_{i} f_{i+} f_{+i}}$$

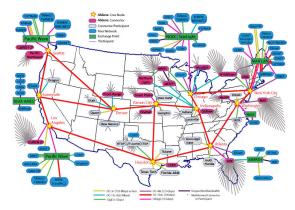
 $\Rightarrow f_{i+}f_{+i} \text{ is the expected fraction of edges joining nodes in } C_i$  $\Rightarrow \text{Random edges preserving degree distribution yields } r_a = 0$ 

• Perfectly assortative mixing yields  $r_a^{\max} = 1$ , while the minimum is

$$r_{a}^{\min} = -\frac{\sum_{i} f_{i+} f_{+i}}{1 - \sum_{i} f_{i+} f_{+i}} > -1$$



- Abilene network for US universities and research labs
  - 'Core' nodes, as well as e.g., 'Connector' nodes and 'Exchange points'



Hierarchical structure, suggestive of disassortative mixing



Tabulated counts of inter-category edges in Abilene

	Core	Exchange	Peer	Conn.	Part.	Conn./Part.
Core	14	6	5	17	0	16
Exchange	6	1	46	2	0	0
Peer	5	46	0	0	0	1
Conn.	17	2	0	0	203	0
Part.	0	0	0	203	0	34
Conn./Part.	16	0	1	34	34	0

▶ Fractions *f<sub>ij</sub>* obtained by scaling table entries by the total of 675

► Assortativity coefficient r<sub>a</sub> = -0.3162, close to r<sub>a</sub><sup>min</sup> = -0.3461 ⇒ Strongly supports our suspicion of disassortative mixing



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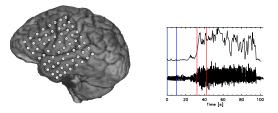
- Epilepsy is the world's most common serious brain disorder
   ⇒ Seizures, i.e., recurrent abnormal neuronal activity
- Ex: Network-oriented analysis of epileptic seizure data in humans
- M. A. Kramer et al, "Emergent network topology at seizure onset in humans," *Epilepsy Res.*, vol. 79, pp. 173-186, 2008
- Leverage few summaries of network characteristics we learnt so far

#### Measurement



Electrode grid (8x8) implanted in the cortical surface of the brain
 ⇒ Also implanted two strips of 6 electrodes (deeper, not shown)

Electrocorticogram (ECoG) data; voltages indicative of brain activity



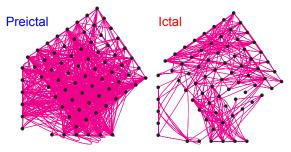
Two 10 sec. intervals of interest for comparison:
 ⇒ Preictal period: prior to the epileptic seizure
 ⇒ lctal period: immediately after start of seizure

► Top time-series is smoothed, averaged over 8 seizure signals

# Network graph construction



- Network → represent couplings among brain regions
   ⇒ Graphs for the preictal and ictal periods, for 8 seizures
   Vertices: correspond to the 76 electrodes (cortical brain regions)
- ► Edges: threshold correlations between pairwise 10 sec. time series



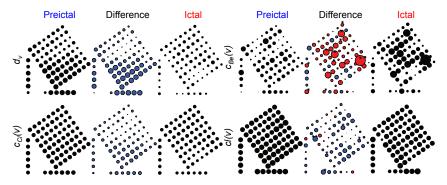
- Brain is in two very different states before and during seizure
  - $\Rightarrow$  Thinning of edges, coupling reduction at seizure onset
  - $\Rightarrow$  Closest to the strips, where seizure was suspected to emanate

### Summaries of network characteristics



Evaluated degree, closeness, betweenness centrality; clustering coeff.

 $\Rightarrow$  Show preictal and ictal periods, as well as their difference



Identifies spatially localized brain regions that may facilitate seizures
 May serve to more precisely guide surgical intervention





- Network cohesion
- Cohesive subgroups
- ► Familiarity
- Reachability
- Robustness
- Local density
- Cliques
- k-plex and k-core
- k-clique and k-club
- Egonet

- Clustering coefficient
- Bridge edges
- Giant connected component
- Small world phenomenon
- Average path length
- Bowtie structure
- Strongly-connected component
- (Dis) assortative mixing
- Homophily
- Brain networks